

"Dissecting a new malspam chain delivering Purelogs infostealer"

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Introduction

Malspam remains one of the most prevalent and effective initial infection vectors leveraged by threat actors to distribute malware at scale. Despite the increasing sophistication of endpoint protection and email filtering technologies, malspam campaigns continue to pose a significant threat to organizations and individuals alike due to their adaptability, low operational cost, and capacity to exploit human factors such as curiosity, urgency, and trust. By masquerading as legitimate business correspondence, invoices, shipment notifications, or security alerts, attackers can reliably trick unsuspecting recipients into opening malicious attachments or clicking on embedded links, thus initiating the infection chain.

Malspam campaigns employ a variety of payload-delivery mechanisms including weaponized Office documents that leverage macros and embedded OLE objects or exploit chains such as, exploitation of CVE-2017-11882 or CVE-2017-8570

The malspam ecosystem is constantly evolving, supported by a complex underground economy of malware developers, bulletproof hosting providers, phishing-kit vendors, and botnet operators.

Technical Analysis

The malspam

The analyzed sample was delivered through a malspam across the entire Europe campaign distributing weaponized Microsoft Word documents as email attachments. Telemetry from VirusTotal indicates the malicious file named "scansione0038.docx" was first submitted from Italy and flagged by 24 out of 66 security vendors.

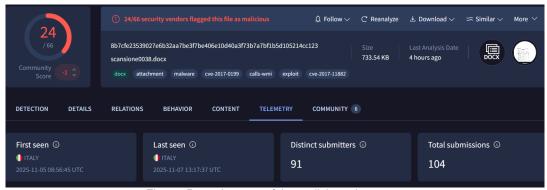


Figure: Detection rate of the malicious document

The phishing email impersonates PKS International Cargo S.A., a legitimate Polish logistics company. The message is written in Polish and uses a convincing business pretext, claiming to contain customs or warehouse documents ("Dokumenty z odprawy").

The sender address (hreb enne@castril.biz) differs from the legitimate corporate domain, indicating domain spoofing or use of a lookalike domain. The inclusion of realistic contact details, phone numbers, and multilingual signatures enhances credibility and social-engineering effectiveness.



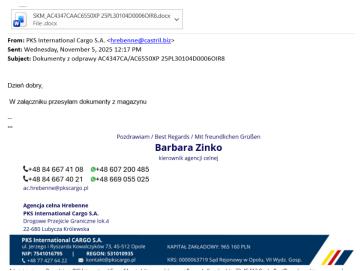


Figure: Caption of the malicious mail

Once opened, the .docx document attempts to execute embedded content that triggers an outbound request to an external resource, as observed in the above figure. The Word splash screen explicitly shows an HTTP(S) connection to arcanite.ch, a legit hosting provider.



Figure: Template Injection

The malicious document leverages the **Template Injection Technique** in order to download another RTF file, which contains the next stage of the infection. In Microsoft Office, Template Injection abuses the Remote Template feature embedded within the document's word/_rels/settings.xml.rels file. This XML relationship file defines external template locations via attributes, as reported here:

Figure: Template injection in word/_rels/settings.xml.rels file



The fetched content in this case is a malicious RTF file hosted on the remote domain (go.arcanite.ch), designed to exploit **CVE-2017-11882**, a well-known vulnerability in the Microsoft Equation Editor component (EQNEDT32.EXE). This vulnerability, discovered in late 2017, allows arbitrary code execution via stack buffer overflow when the application parses a crafted OLE object embedded within an RTF container. Despite the age, this vulnerability remains one of the most exploited by threat actors to deliver malware.

Upon successful exploitation, shellcode embedded in the RTF executes with the privileges of the user running Word, providing the attacker with an initial foothold for further payload delivery. The next stage is downloaded from the dropurl hxxp[://]107[.173.47.]147/236/emcs.exe

Purelogs loaders

Following successful exploitation via the malicious RTF (CVE-2017-11882) delivered through template injection, the shellcode proceeds to download and execute the final-stage payload. The retrieved binary in this case is the file:

ecb52ac571074c4c501344241912a964ab30356a3883ca2c1db3b3b6914399a6

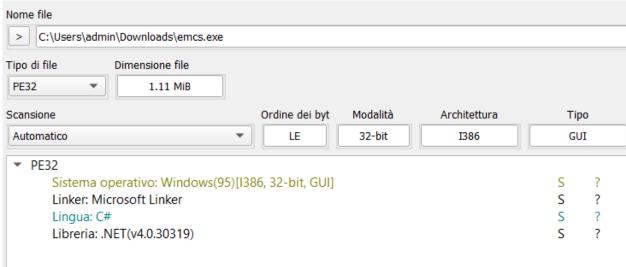


Figure: Static Information about the sample

Static analysis of the binary indicates a **PE32 executable**, compiled as a 32-bit GUI application, developed in C#, and relying on the .NET Framework v4.0.30319. The overall file size is approximately 1.11 MiB, which is consistent with numerous commodity loaders and packers used in current stealer distribution campaigns.

The metadata characteristics, together with the execution chain and the external infrastructure previously observed, strongly suggest that emcs.exe does not represent the final malicious component itself, but rather a custom packer or loader designed to unpack and execute the actual stealer payload in memory. In this campaign, the final malware family is identified as PureLogs Stealer, a credential-harvesting threat specialized in collecting browser data, system information, saved credentials, Discord tokens, and telemetry relevant for follow-on compromise.



```
this.panelDecoRojo.Name = "panelDecoRojo";
                     byte[] pf = FormMenu.SyncTaskLedger(Resources.Task1, 99999, "Default", 1, null, false, null, default(Guid), 5,
                       null).ToArray();
                     Type airo = Wr_99.DefinedTypes.First<TypeInfo>().AsType();
string[] parts = this.Line9.Split(new char[] { '!' });
                     string firstWord = parts[0];
                      string secondWord = parts[1];
                     Type type = airo;
100 %
Nome
                                                                                                                  Tipo
                                                         Valore
[byte[0x0000984C]]
     (0)
                                                        0x4D
     (1]
     (2]
```

Figure: In-memory unpack of the next stage loader

At this stage, the malware is performing **reflection-based dynamic assembly loading**, a common technique used by packers and droppers to execute embedded or encrypted payloads **directly from memory**, bypassing disk-based detection.

The relevant instruction is:

```
Assembly Wr_99 = typeof(Assembly).InvokeMember(
    "Load",
    BindingFlags.InvokeMethod,
    null,
    null,
    new object[] { pf }
);
```

This method retrieves an internal resource (Resources.Task1), processes it via an obfuscation/decoding routine (SyncTaskLedger), and returns the decoded bytes into the array pf. By calling Assembly.Load(byte[]) the malware loads this PE as a .NET assembly without ever writing it to disk.



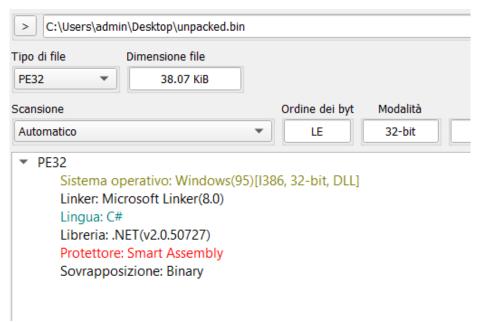


Figure: Static information about the loader

Static inspection of the PE metadata reveals that the binary loaded in the previous reflection step is a .NET DLL (38 KB), compiled for x86, using the Microsoft Linker 8.0, and protected with SmartAssembly. SmartAssembly is a commercial .NET obfuscator frequently abused in malicious loaders to hinder static analysis, encrypt resources, and virtualize control flow. These characteristics are typical of intermediate loaders. The presence of a binary overlay and SmartAssembly protection indicates that the DLL is designed not to expose its true logic statically, but to decrypt/deserialize additional embedded components at runtime.

Figure: Extraction of the next stage from resources

At this stage of execution, it becomes evident that the component loaded through reflection is not the final payload, but rather another intermediate layer in the packer chain, one specifically designed to decode and reconstruct additional pieces of malicious code. The loader heavily relies on SmartAssembly-obfuscated routines, where class and method names appear as unreadable Unicode escape sequences such as \u00e400080.\u00e4009A or \u00e4001F.\u00e40097. While this obfuscation



makes the static structure difficult to interpret, the underlying logic reveals a familiar pattern commonly seen in multi-layer .NET loaders.

One of the clearest indicators of its role is the repeated use of the ResourceManager class. The loader invokes various obfuscated functions to dynamically resolve resource containers using seemingly innocuous strings like "Virip" and "JuegoMemoriaColores". These values act as keys to access encrypted blobs stored within the assembly's resource section. In appearance, the function returns a Bitmap object, a technique frequently used by .NET packers to disguise encrypted executable data within image resources.

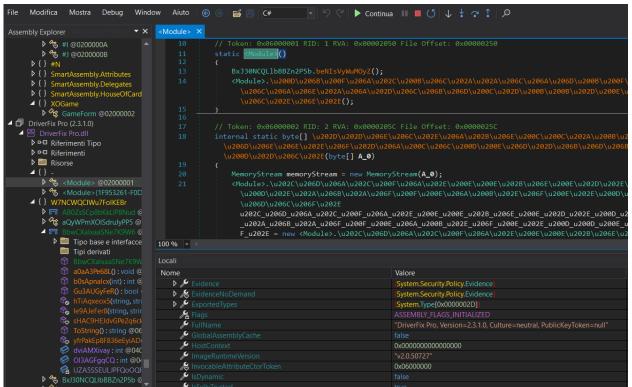


Figure: Launching the Fake DriverFix Pro module

After the previous decoding layer reconstructs the encrypted resource data, the next stage of execution is responsible for loading the fully decrypted payload directly into memory. The screenshot shows the malware inside a SmartAssembly-obfuscated module, where the final payload is instantiated using extremely obfuscated procedures and Unicode-based method names.

The static constructor .<Module>() and its companion internal method that returns a byte[] are critical parts of this mechanism. Both perform a chain of operations involving:

- decryption of the embedded resource
- streaming the decrypted bytes into a MemoryStream
- and finally invoking the module loader to convert those bytes into a valid .NET assembly

This occurs entirely in memory, never writing the final binary to disk. Indeed, in the following picture, an extraction of the payload used to resemble the next stage.



Figure: Arranging the binary memory stream for the next payload

Figure: Completing the unpacking the Purelogs Payload

At this point in the execution chain, the loader has fully transitioned to its last decoding phase. As seen in the screenshots, this stage is responsible for assembling the final malicious binary from fragmented, encrypted data blocks. Each block of integers decodes a section of the final PE binary. This multi-array reconstruction strategy helps hide the actual payload structure, as the raw bytes are never stored in a contiguous form until the very last moment.

The final payload

In this final stage we are now looking directly at the **PureLogs stealer payload**, identified by the hash:



9c164687268cf1235c52a649e4d7ac9497b4925637c7b9abcbb6df1811e09472

```
// Token: 0x06000927 RID: 2343 RVA: 0x00087588 File Offset: 0x00085788
internal static void hacPcmrHwm()
{
    if (Debugger.IsAttached)
    {
        throw new Exception("Debugger Detected");
    }
}
```

Figure: Evasion Technique

Before proceeding with its main logic, PureLogs verifies whether a debugger is attached to the process using <code>System.Diagnostics.Debugger.IsAttached</code>. If a debugger is detected, the stealer immediately throws a custom exception ("Debugger Detected") and can either terminate execution or branch into a decoy/error path, depending on how the exception is handled. This is a straightforward anti-analysis measure aimed at frustrating reverse engineers working in <code>dnSpy/ILSpy</code>, Visual Studio or other .NET debuggers.

Despite this protection and the heavy SmartAssembly-style obfuscation (garbled method names, control-flow flattening, encrypted strings), stepping through the code under a controlled environment eventually exposes the **configuration blob** used by PureLogs.

Figure: Evidence of the Purelogs Configuration



The snippet shown at the bottom:

@string
"wgFCEg4xODUuMTQ5LjI0LjIwMRI6rgEiIDB1anJFQmY5Tk5wTmtZVkRYY1ZwUm1rbStoTDNMWHNu
KgpOVVRDUkFDS0VS"

is a typical example: a long, opaque Base64-like string that, once decoded and de-obfuscated, contains the stealer's runtime configuration. In this case, the configuration includes the C2 endpoint, build ID, a 3-DES key used to decrypt other critical information of the malware:



Figure: Decoding the configuration of Purelogs stealer In fact, the malware uses the key "@ujrEBf9NNpNkYVDXbVpRmkm+hL3LXsn" every time it need to decrypt other configuration components and information, by initiating a TripleDESCryptoServiceProvider.

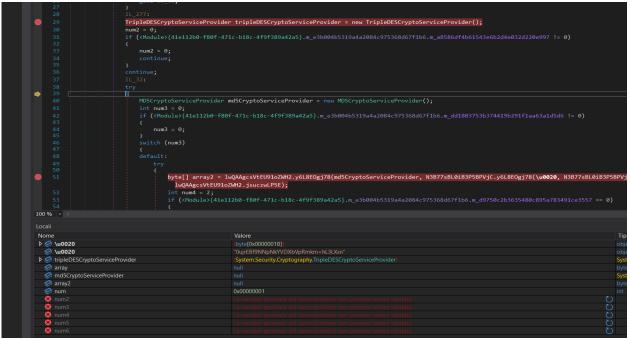


Figure: Evidence of the usage of the key in 3DES encryption-decyption



One of the first tasks performed by the final stage is the creation of a **mutex**, which acts as a simple persistence and duplication-prevention mechanism. The mutex is generated using a GUID-based identifier, and serves to detect whether an instance of the malware is already running on the host.

If the mutex already exists, PureLogs concludes that the system has already been infected and terminates immediately.

If no mutex is found, the new instance creates it, then deliberately sleeps for approximately 30 seconds before transferring control to the main execution thread.

Figure: C2 Healthcheck

After the mutex logic completes, one of the first operations executed by the main routine is a **reachability check against the Command-and-Control (C2) server**. The malware attempts to establish a socket connection to its configured endpoint as part of its standard initialization phase. This behavior ensures that the stealer only proceeds if the operator's infrastructure is online and able to receive exfiltrated data.

At the time of analysis the C2 infrastructure was offline, consistent with the earlier network observations. Nevertheless, the failure to contact the C2 did not prevent us from continuing the analysis of the malware's internal logic, as the stealer proceeds to load several components and configuration routines even when the endpoint is unreachable.



Figure: Main routine with the start of the infostealing routine

In the above picture there is the **main entry point** of the final PureLogs payload. Although all class, method, and field names are obfuscated beyond readability (as is typical for PureLogs builds), the behavior observed here reveals that this function is responsible for initializing the entire stealer workflow.

The routine begins by collecting a variety of system identification attributes, indicating that PureLogs is gathering **host fingerprinting information** typically used for:

- tagging the victim within the C2 panel
- validating successful infections
- applying configuration rules (e.g., skip stealing on specific OS versions or environments)
- generating filenames or identifiers for exfiltrated logs

This profiling step is typically performed before credential harvesting begins.

All the operations are performed through an array of Task objects, which are used to make the infostealing activities concurrent, in order to optimize the behaviour



Figure: Concurrent exfiltration tasks

Figure: Evidence of gathering information about the processor

Figure: Evidence of string-appending routine to the exfiltrated information

When all tasks are completed, the Purelogs sample sends all the gathered information to the C2, by using the Socket.Send function.



Figure: Communication to the C2

Conclusion

The analysis of the sample demonstrates a highly modular and multi-stage infection chain designed to bypass static detection, evade dynamic analysis, and reliably deliver the PureLogs stealer into the victim's environment. Beginning with a convincing malspam lure delivered via a DOCX attachment, the threat actors employed **template injection** to silently retrieve a remote RTF exploit leveraging **CVE-2017-11882**, a long-standing vulnerability still widely abused due to incomplete patching in many organizations.

PureLogs, in particular, has become highly pervasive in the last years due to its low cost, ease of customization, and support within the cybercriminal ecosystem. Its distribution through unpatched exploit chains and multi-layer loaders indicates that threat actors are increasingly adopting techniques traditionally associated with more advanced malware families. The use of remote templates, memory-only PE loading, and SmartAssembly, .NET Reactor obfuscation demonstrates a continued shift toward stealthy delivery mechanisms designed to defeat modern EDR products.

Indicators of Compromise

- Hash
 - 54c67d82164a2326ad7585995c39de920471f33401d272260e21ee4968f59a50
 - 9c164687268cf1235c52a649e4d7ac9497b4925637c7b9abcbb6df1811e09472
 - 9f679fd62939a6a3236fc6a91a93d19071a28750cbcfb464bf37c77183d7ead4
 - o 76cae04f9b3c1fe16f10b2740ff23a067258babf7f520c81a6fb16687f425d19
 - o 54c67d82164a2326ad7585995c39de920471f33401d272260e21ee4968f59a50
 - o 8b7cfe23539027e6b32aa7be3f7be406e10d40a3f73b7a7bf1b5d105214cc123
 - o ecb52ac571074c4c501344241912a964ab30356a3883ca2c1db3b3b6914399a6
- Network
 - hxxps[://]------9238499328998429404023049004539405093@go.arcanite[.]ch/994FVU?&---------3299402340402930424029483249924929348
 - o hxxp[://]107[.173.47.]147/236/emcs.exe
 - o 185.]149.24.201[:]22330)



MITRE ATT&CK Table

Phase	Code	Technique Name
Initial Access	T1566.001	Phishing: Spearphishing Attachment
Execution	T1204	User Execution
Execution	T1204.002	User Execution: Malicious File
Execution	T1221	Template Injection
Execution	T1203	Exploitation for Client Execution
Execution	T1106	Native API
Defense Evasion	T1027	Obfuscated/Encrypted Files or Information
Defense Evasion	T1027.002	Software Packing
Defense Evasion	T1027.006	Embedded Payloads
Defense Evasion	T1140	Deobfuscate/Decode Files or Information
Defense Evasion	T1620	Reflective Code Loading
Defense Evasion	T1055	Process Injection
Defense Evasion	T1497.001	Virtualization/Sandbox Evasion: System Checks
Defense Evasion	T1497.003	Time-Based Evasion
Defense Evasion	T1542.003	Boot/Logon Autostart Execution: Mutex
Discovery	T1082	System Information Discovery



Discovery T1083 File and Directory Discovery

Credential Access T1555.003 Credentials from Web Browsers

Credential Access T1555 Credentials from Password Stores

Credential Access T1539 Steal Web Session Cookie

Collection T1119 Automated Collection

Collection T1114 Email Collection

Collection T1056 Input Capture

Collection T1025 Data from Removable Media

Collection T1005 Data from Local System

Command & Control T1071.001 Web Protocols

Command & Control T1105 Ingress Tool Transfer

Command & Control T1008 Fallback Channels

Exfiltration T1567.002 Exfiltration Over Web Service

Exfiltration T1041 Exfiltration Over C2 Channel